MIPS Architecture and Assembly Language Overview

Adapted from: http://edge.mcs.dre.g.el.edu/GICL/people/sevy/architecture/MIPSRef(SPIM).html

[Register Description] [I/O Description]

Data Types and Literals

Data types:

- Instructions are all 32 bits
- byte(8 bits), halfword (2 bytes), word (4 bytes)
- a character requires 1 byte of storage
- an integer requires 1 word (4 bytes) of storage

Literals:

- numbers entered as is. e.g. 4
- characters enclosed in single quotes. e.g. 'b'
- strings enclosed in double quotes. e.g. "A string"

Registers

- 32 general-purpose registers
- register preceded by \$ in assembly language instruction two formats for addressing:
 - using register number <u>e.g.</u> \$0 through \$31
 - using equivalent names <u>e.g.</u> \$t1, \$sp
 - special registers Lo and Hi used to store result of multiplication and division
 - not directly addressable; contents accessed with special instruction mfhi ("move from Hi") and mflo ("move from Lo")
- · stack grows from high memory to low memory

This is from Figure 9.9 in the Goodman&Miller text

Register Number	Alternative Name	Description			
0	zero	the value 0			
1	\$at	(assembler temporary) reserved by the assembler			
2-3	\$v0 - \$v1	(values) from expression evaluation and function results			
4-7	\$a0 - \$a3	(arguments) First four parameters for subroutine. Not preserved across procedure calls			
8-15	\$t0 - \$t7	(temporaries) Caller saved if needed. Subroutines can use w/out saving. Not preserved across procedure calls			
16-23	\$s0 - \$s7	(saved values) - Callee saved. A subroutine using one of these must save original and restore it before exiting. Preserved across procedure calls			
24-25	\$t8 - \$t9	(temporaries) Caller saved if needed. Subroutines can use w/out saving. These are in addition to \$t0 - \$t7 above. Not preserved across procedure calls.			
26-27	\$k0 - \$k1	reserved for use by the interrupt/trap handler			
28	\$gp	global pointer. Points to the middle of the 64K block of memory in the static data segment.			
29	\$sp	stack pointer Points to last location on the stack.			
30	\$s8/\$fp	saved value / frame pointer Preserved across procedure calls			
31	\$ra	return address			

See also Britton section 1.9, Sweetman section 2.21, Larus Appendix section A.6

Program Structure

- just plain text file with data declarations, program code (name of file should end in suffix .s to be used with SPIM simulator)
- data declaration section followed by program code section

Data Declarations

- placed in section of program identified with assembler directive .data
- declares variable names used in program; storage allocated in main memory (RAM)

Code

• placed in section of text identified with assembler directive .text

- contains program code (instructions)
- starting point for code e.g.ecution given label main:
- ending point of main code should use exit system call (see below under System Calls)

Comments

- anything following # on a line
 - # This stuff would be considered a comment
- Template for a MIPS assembly language program:
 - # Comment giving name of program and description of function
 # Template.s

#	Bare-bones	outline	of	MIPS	assembly	language	program	

.data	<pre># variable declarations follow this line #</pre>	
.text	<pre># instructions follow this line</pre>	
main:	<pre># indicates start of code (first instruction to execut #</pre>	e)
<pre># End of program, leav</pre>	a blank line afterwards to make SPIM happy	

Data Declarations

format for declarations:

name: storage_type value(s)

- create storage for variable of specified type with given name and specified value
- value(s) usually gives initial value(s); for storage type .space, gives number of spaces to be allocated

Note: labels always followed by colon (:)

example			
varl: arrayl:	.word .byte		<pre># create a single integer variable with initial value 3 # create a 2-element character array with elements initialized</pre>
-	-		# to a and b
array2:	.space	40	<pre># allocate 40 consecutive bytes, with storage uninitialized # could be used as a 40-element character array, or a # 10-element integer array; a comment should indicate which!</pre>

Load / Store Instructions

- · RAM access only allowed with load and store instructions
- all other instructions use register operands

load:

lw register destination, RAM source

#copy word (4 bytes) at source RAM location to destination register.

lb register_destination, RAM_source

#copy byte at source RAM location to low-order byte of destination register, # and sign-e.g.tend to higher-order bytes

store word:

sw register source, RAM destination

#store word in source register into RAM destination

sb register_source, RAM_destination

#store byte (low-order) in source register into RAM destination

load immediate:

li register_destination, value

#load immediate value into destination register

example	:		
var1:	.data .word	23	# declare storage for varl; initial value is 23
	.text		
start	:		
	lw	\$t0, var1	<pre># load contents of RAM location into register \$t0: \$t0 = var1</pre>
	li	\$t1, 5	<pre># \$t1 = 5 ("load immediate")</pre>
	SW	\$t1, var1	<pre># store contents of register \$t1 into RAM: var1 = \$t1</pre>
	done		

Indirect and Based Addressing

· Used only with load and store instructions

```
load address:
```

la \$t0, var1

 copy RAM address of var1 (presumably a label defined in the program) into register \$10

```
indirect addressing:
```

lw \$t2, (\$t0)

• load word at RAM address contained in \$t0 into \$t2

sw \$t2, (\$t0)

• store word in register \$t2 into RAM at address contained in \$t0

based or indexed addressing:

lw \$t2, 4(\$t0)

- load word at RAM address (\$t0+4) into register \$t2
- "4" gives offset from address in register \$t0

sw \$t2, -12(\$t0)

- store word in register \$t2 into RAM at address (\$t0 12)
- negative offsets are fine

Note: based addressing is especially useful for:

- arrays; access elements as offset from base address
- · stacks; easy to access elements at offset from stack pointer or frame pointer

example

```
.data
                        12
array1:
                .space
                                         # declare 12 bytes of storage to hold array of 3 integers
                 .text
start:
                la
                         $t0, array1
                                                  #
                                                    load base address of array into register $t0
                         $t1, 5
                li
                                         #
                                            $t1 = 5
                                                      ("load immediate"
                sw $t1, ($t0)
li $t1, 13
                                            first array element set to 5; indirect addressing
                                         #
                                             $t1 = 13
                                         #
                sw $t1, 4($t0)
                                         #
                                            second array element set to 13
                li $t1, -7
                                         #
                                             st1 = -7
                sw $t1, 8($t0)
                                         #
                                            third array element set to -7
                done
```

Arithmetic Instructions

- · most use 3 operands
- all operands are registers; no RAM or indirect addressing
- operand size is word (4 bytes)

```
add
            $t0,$t1,$t2
                                      #
                                         $t0 = $t1 + $t2;
                                                                         add as signed (2's complement) integers
                                          $t2 = $t3 Đ $t4
                                      #
sub
            $t2,$t3,$t4
addi
            $t2,$t3, 5
                                      # $t2 = $t3 + 5;
                                                                      "add immediate" (no sub immediate)
            $t1,$t6,$t7
addu
                                      #
                                         $t1 = $t6 + $t7;
                                                                         add as unsigned integers
                                          $t1 = $t6 + $t7;
                                                                         subtract as unsigned integers
subu
            $t1,$t6,$t7
                                      #
                                          multiply 32-bit quantities in $t3 and $t4, and store 64-bit
result in special registers Lo and Hi: (Hi,Lo) = $t3 * $t4
Lo = $t5 / $t6 (integer quotient)
Hi = $t5 mod $t6 (remainder)
move quantity in special register Hi to $t0: $t0 = Hi
move quantity in special register Lo to $t1: $t1 = Lo
used to get at regult of product or gutient
mult
            $t3,$t4
                                      #
            $t5,$t6
div
                                      #
                                      #
mfhi
            $t0
                                      #
mflo
            $t1
                                      #
                                          used to get at result of product or quotient
            $t2,$t3 # $t2 = $t3
move
```

Control Structures

Branches

· comparison for conditional branches is built into instruction

btarget#unconditional branch to program label targetbeq\$t0,\$t1,target#branch to target if\$t0 = \$t1blt\$t0,\$t1,target#branch to target if\$t0 < \$t1</td>ble\$t0,\$t1,target#branch to target if\$t0 <= \$t1</td>bgt\$t0,\$t1,target#branch to target if\$t0 <= \$t1</td>

```
branch to target if
                                                $t0 >= $t1
bge
        $t0,$t1,target
                        #
bne
        $t0,$t1,target
                        #
                          branch to target if $t0 <> $t1
```

Jumps

unconditional jump to program label target target # jump to address contained in \$t3 ("jump register") \$t3

Subroutine Calls

subroutine call: "jump and link" instruction

i jr

- ial sub label # "jump and link"
- copy program counter (return address) to register \$ra (return address register)
- jump to program statement at sub_label

subroutine return: "jump register" instruction

"jump register" \$ra jr

• jump to return address in \$ra (stored by jal instruction)

Note: return address stored in register \$ra; if subroutine will call other subroutines, or is recursive, return address should be copied from \$ra onto stack to preserve it, since jal always places return address in this register and hence will overwrite previous value

System Calls and I/O (SPIM Simulator)

· used to read or print values or strings from input/output window, and indicate program end

- use syscall operating system routine call
- first supply appropriate values in registers \$v0 and \$a0-\$a1

• result value (if any) returned in register \$v0

The following table lists the possible syscall services.

Service	Code in \$v0	Arguments	Results
print_int	1	\$a0 = integer to be printed	
print_float	2	\$f12 = float to be printed	
print_double	3	\$f12 = double to be printed	
print_string	4	\$a0 = address of string in memory	
read_int	5		integer returned in \$v0
read_float	6		float returned in \$v0
read_double	7		double returned in \$v0
read_string	8	\$a0 = memory address of string input buffer \$a1 = length of string buffer (n)	
sbrk	9	\$a0 = amount	address in \$v0
exit	10		

• The print_string service expects the address to start a null-terminated character string. The directive .asciiz creates a null-terminated character string.

- 0 The read_int, read_float and read_double services read an entire line of input up to and including the newline character. 0
 - The read_string service has the same semantices as the UNIX library routine fgets. • It reads up to n-1 characters into a buffer and terminates the string with a null character.
 - If fewer than n-1 characters are in the current line, it reads up to and including the newline and terminates the string with a null character.
- The sbrk service returns the address to a block of memory containing n additional bytes. This would be used for dynamic memory allocation.

• The exit service stops a program from running.

.text

\$v0, 4

1i

e.g.

string1

main:

F

Print out integer value contained in register \$t2 e.a.

li	\$v0,	1	#	load	appropriate system call code into register \$v0;
			#	code	for printing integer is 1
move	\$a0,	\$t2	#	move	integer to be printed into \$a0: \$a0 = \$t2
syscall			#	call	operating system to perform operation

e.g. Read integer value, store in RAM location with label int_value (presumably declared in data section)

	li	\$v0, 5	<pre># load appropriate system call code into register \$v0; # code for reading integer is 5</pre>
	syscall # ca sw \$v0, int_value # va	<pre># call operating system to perform operation # value read from keyboard returned in register \$v0; # store this in desired location</pre>	
Print ou	t string	(useful for prompts)	
	.data .asciiz	"Print this.\n"	<pre># declaration for string variable, # .asciiz directive makes string null terminated</pre>

load appropriate system call code into register \$v0;

- # code for printing string is 4la\$a0, string1# load address of string to be printed into \$a0syscall# call operating system to perform print operation

e.g. To indicate end of program, use **exit** system call; thus last lines of program should be:

li \$v0,10 # system call code for exit = 10 syscall # call operating sys